## 1 Non-crossing cycles and Catalan number

This time we will show the following lemma:

**Lemma 1.1.** The number of non-crossing cycles of length k in  $\{1, \ldots, n\}$  is equal to  $C_{k/2}n(n-1)\cdots(n-k/2)$ , where  $C_{k/2}$  is the number of Dyck words of length k.

We will use a slightly different formulation, and will not follow [1] completely. How to think of non-crossing cycles? We observe that a cycle of length k is non-crossing if and only if there exists a tree in the complete graph of order n with k/2 edges and k/2 + 1 vertices, such that the cycle lies in the tree and traverses each edge in the tree exactly twice. It will be clearer if I draw a picture, but I am too lazy to do that here.

Why such a cycle is called non-crossing? It is because one can draw the cycles as follows. Let  $(i_1, \ldots, i_k)$  be a cycle of length k and arrange  $1, \ldots, k$  on the circle. Draw a line segment between a and b if  $i_a = i_b$ . If there is no crossing (and the number of line segment is k/2-1), the cycle will be non-crossing. Again it will be clear if one draws a picture.

Instead of Dyck words, we will consider Dyck paths. A Dyck path of length k (where k is even) is a diagonal lattice path from (0,0) to (k,0) that never goes below the x-axis. Note that the number of rooted trees with k/2 + 1 vertices is equal to the number of Dyck paths with k steps. To see this, we can construct a tree from the Dyck path as follows. If the path goes up, then we make a new branch from the vertex we are staying at, and if it goes down, then we just return to the parent of the vertex.

Since a non-crossing cycle gives rise to an ordered sequence and a Dyck path, and there is an obvious bijection between Dyck paths and Dyck words, this proves Lemma 1.1.

We have shown last time that for a Hermitian Wigner matrix M with entries having mean 0, variance 1, and bounded above by  $o(\sqrt{n})$  in magnitude, then  $\operatorname{Etr}(M^k) = (C_{k/2} + o_k(1))n^{k/2+1}$  (I think in the seminar I said that this seems to be just the upper bound, but indeed this is a 2-sided bound). Combining with  $||M||^k \leq \operatorname{tr}(M^k) \leq n||M||^k$ , we have

$$\mathbf{E} \|M\|^{k} \ge (C_{k/2} + o_{k}(1))n^{k/2}.$$

How to find  $C_{k/2}$ ? We can easily prove by induction that there exists a recurrence relation

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}.$$

Define the generating function of  $C_n$ :

$$f(z) = \sum_{n=0}^{\infty} C_n z^n.$$

Then

$$f(z)^{2} = \sum_{n=0}^{\infty} \left( \sum_{i=0}^{n} C_{i} C_{n-i} \right) z^{n} = \sum_{n=0}^{\infty} C_{n+1} z^{n}.$$

Thus we have the functional equation

$$z(f(z))^2 = f(z) - 1.$$

Solving for f, we have

$$f(z) = \frac{1 - \sqrt{1 - 4z}}{2z},$$

and by finding the coefficients of the Taylor series of f we see that

$$C_n = \frac{1}{n+1} \binom{2n}{n}.$$

## References

- Tao, Terence. Topics in random matrix theory. Graduate Studies in Mathematics, 132. American Mathematical Society, Providence, RI, 2012. x+282 pp. ISBN: 978-0-8218-7430-1
- [2] Thompson, Brady. Talk on 2/23.